

CLEARSY Safety Platform



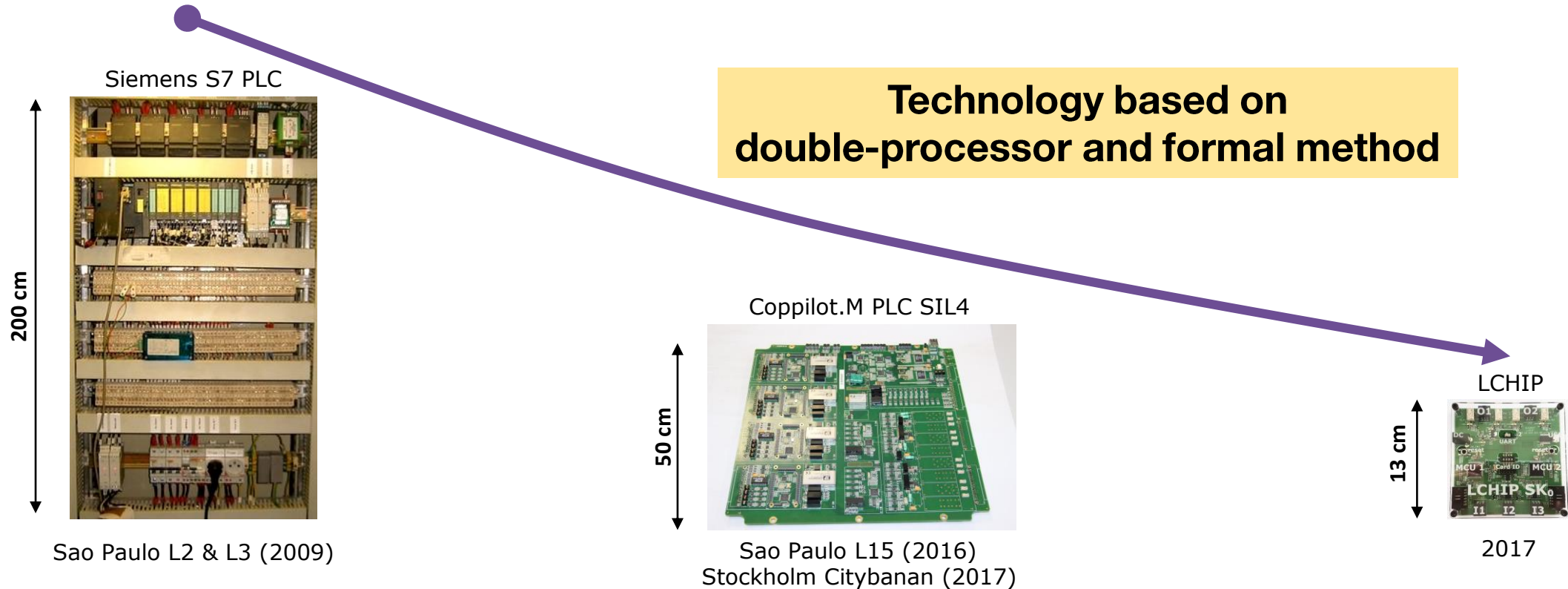
Niteroi, May 10th 2018

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Genesis

- Need for a technical solution to overcome difficulties to develop SIL3/SIL4 systems
 - Require rare human resources to complete successfully
 - Very short delays (~6 months) to design new systems
 - Off-the-shelf block solutions difficult to adapt



Genesis

Technology experimented several times



Platform screen doors controller installed in Stockholm (Citybanan)



Sao Paulo L15 (2016)



Stockholm Citybanan (2017)



Gateway SATURN (2016)

Technology certifiable



COPPILOT.M Stockholm application « série A »
implementing the SIL3 safety function
"Automatic Sliding Doors (ASD) Opening Authorization"

implementing the SIL3 safety function
Certificate N°: 6393741
Date of issue: 03rd March, 2017



Genesis

Next Step

Develop a generic version of this technology that could

- fit a broader range of applications
- be used for educational purpose

CLEARSY Safety Platform in a Nutshell

SOFTWARE & HARDWARE
PLATFORM
FOR **SAFETY CRITICAL** APPLICATIONS

Factory
to generate a function
that executes safely

DESIGN

EXECUTION

≡ ease the development of safety critical applications

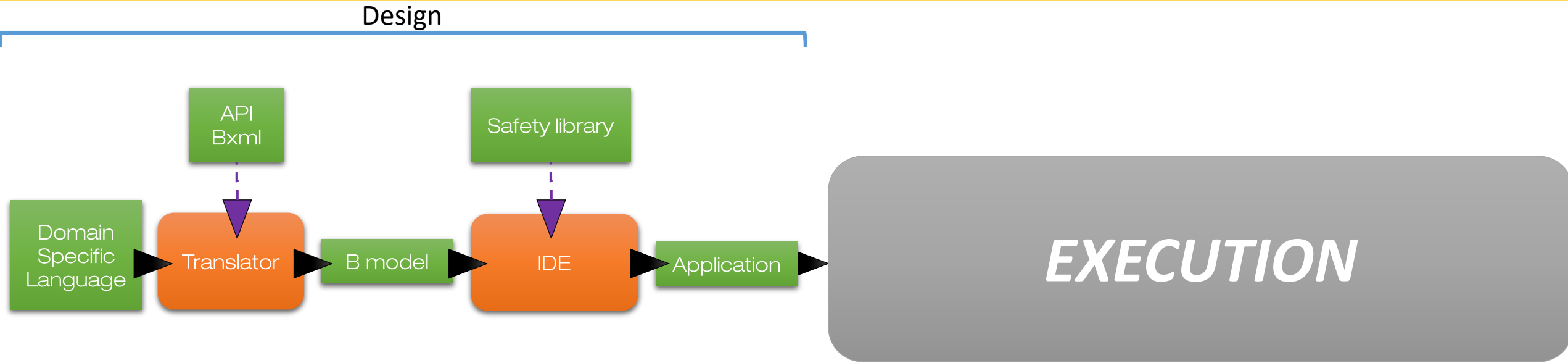
≡ ease the certification of safety critical applications

- Cover the whole development cycle
- Safety principles are built-in
- Based on a formal language and related proof tools
- Mathematical proof replaces unit and integration testing
- Avoid expert transactions over multiple languages

- Safety cannot be altered by the developer
- Comes along with a certification kit

Disclaimer: The developer is responsible for the safety demonstration

CLEARSY Safety Platform in a Nutshell

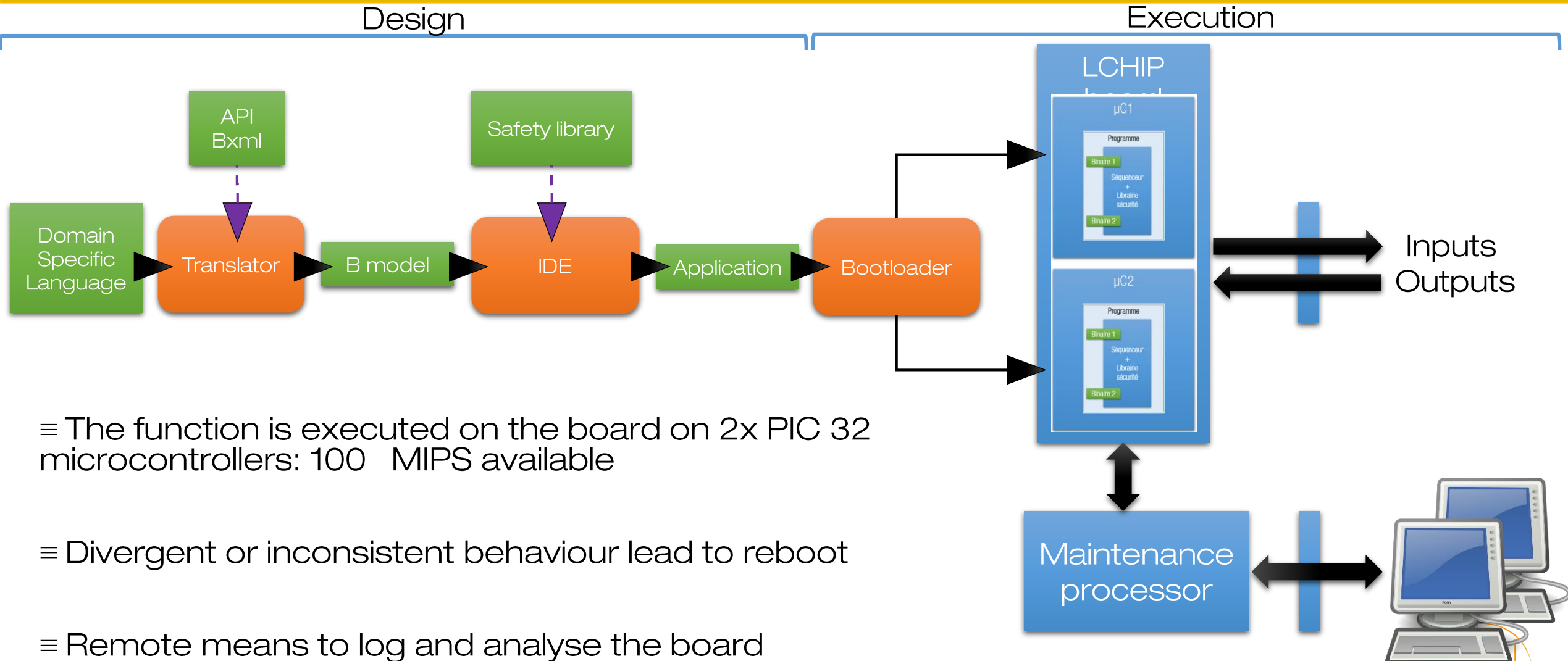


≡ Developer focus on the function to develop by using its usual tools / language

≡ Transformation of the DSL inputs into B formal models, proof, code generation and compilation are fully automatic

Some constraints to comply though

CLEARSY Safety Platform in a Nutshell



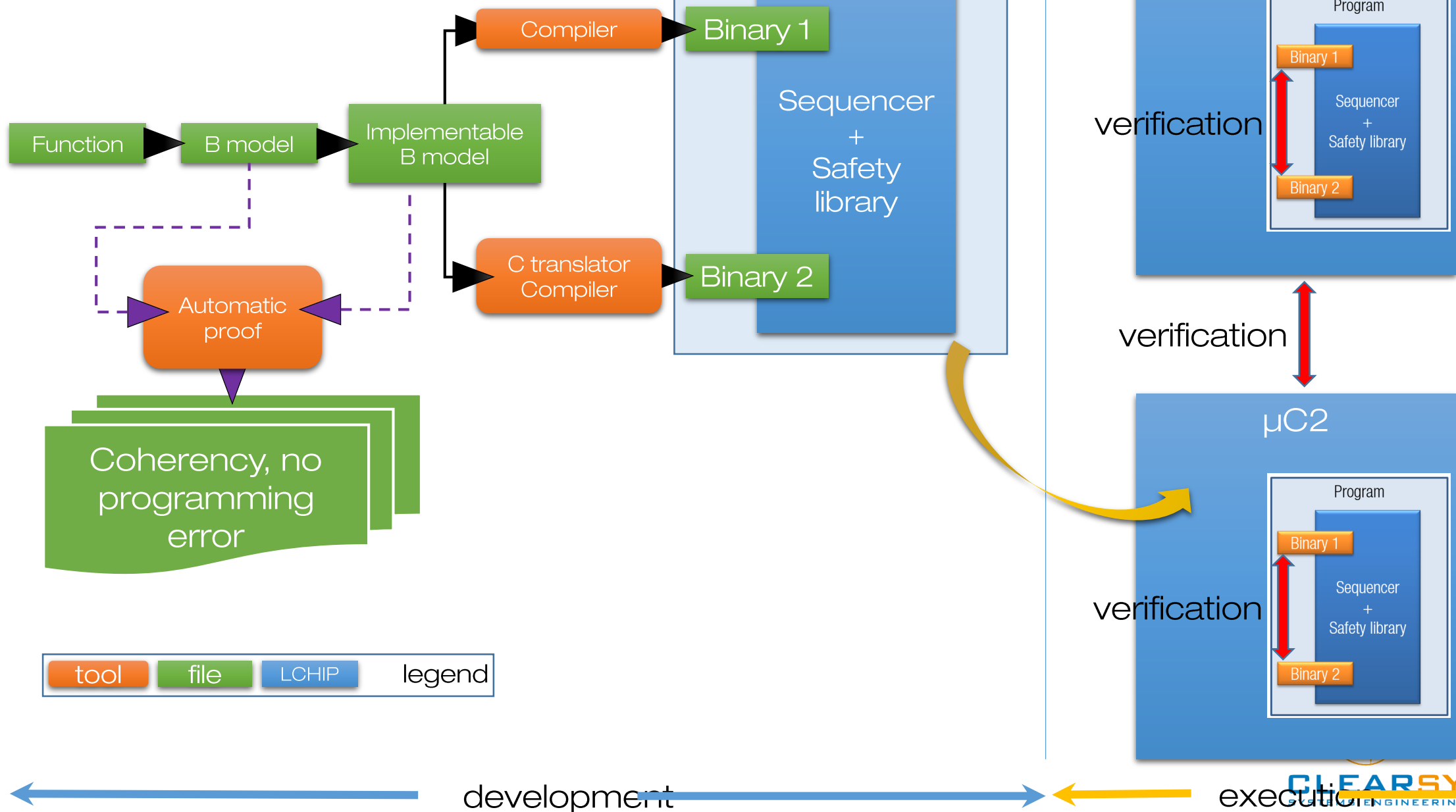
≡ The function is executed on the board on 2x PIC 32 microcontrollers: 100 MIPS available

≡ Divergent or inconsistent behaviour lead to reboot

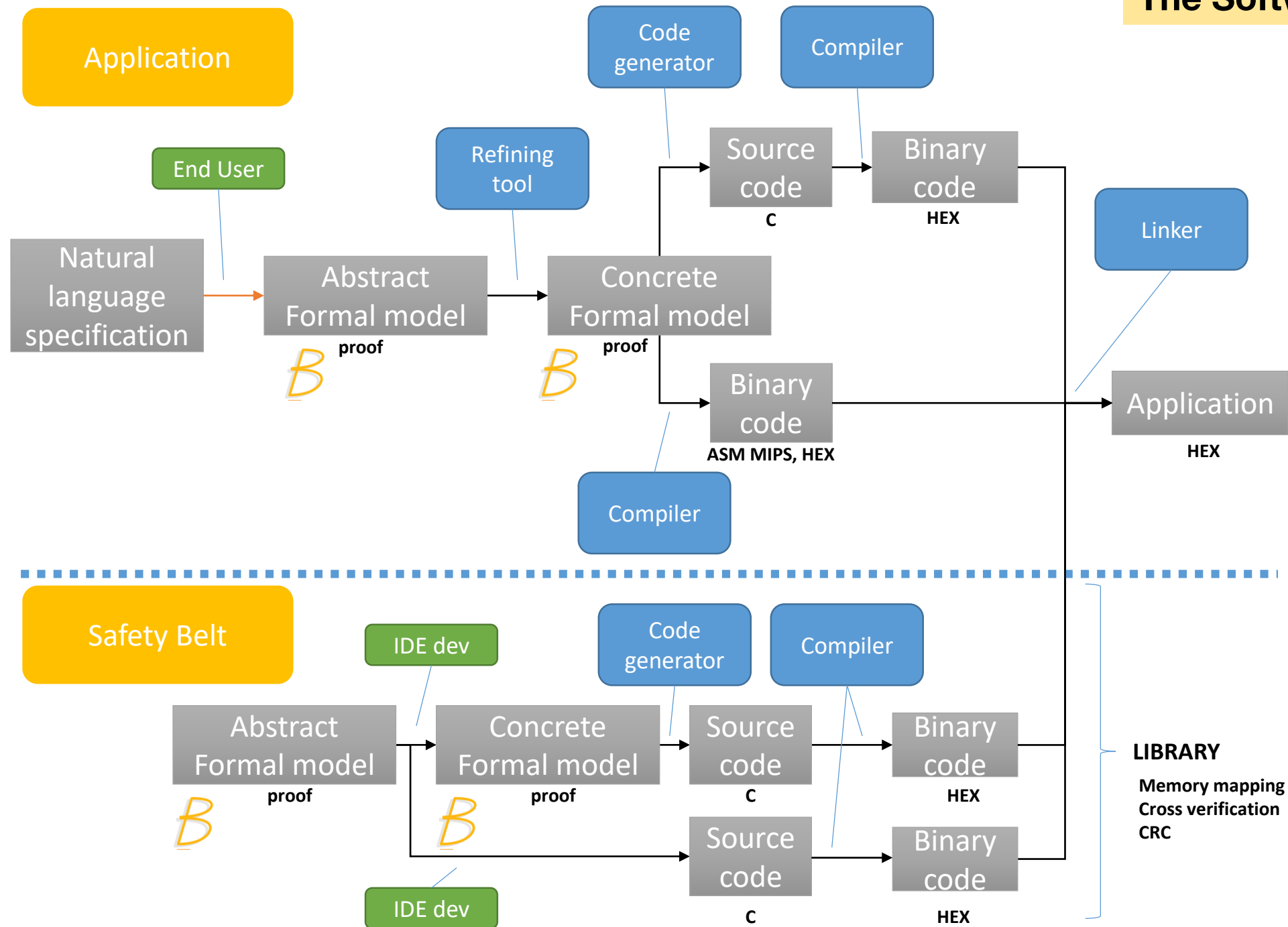
≡ Remote means to log and analyse the board

Technical principles

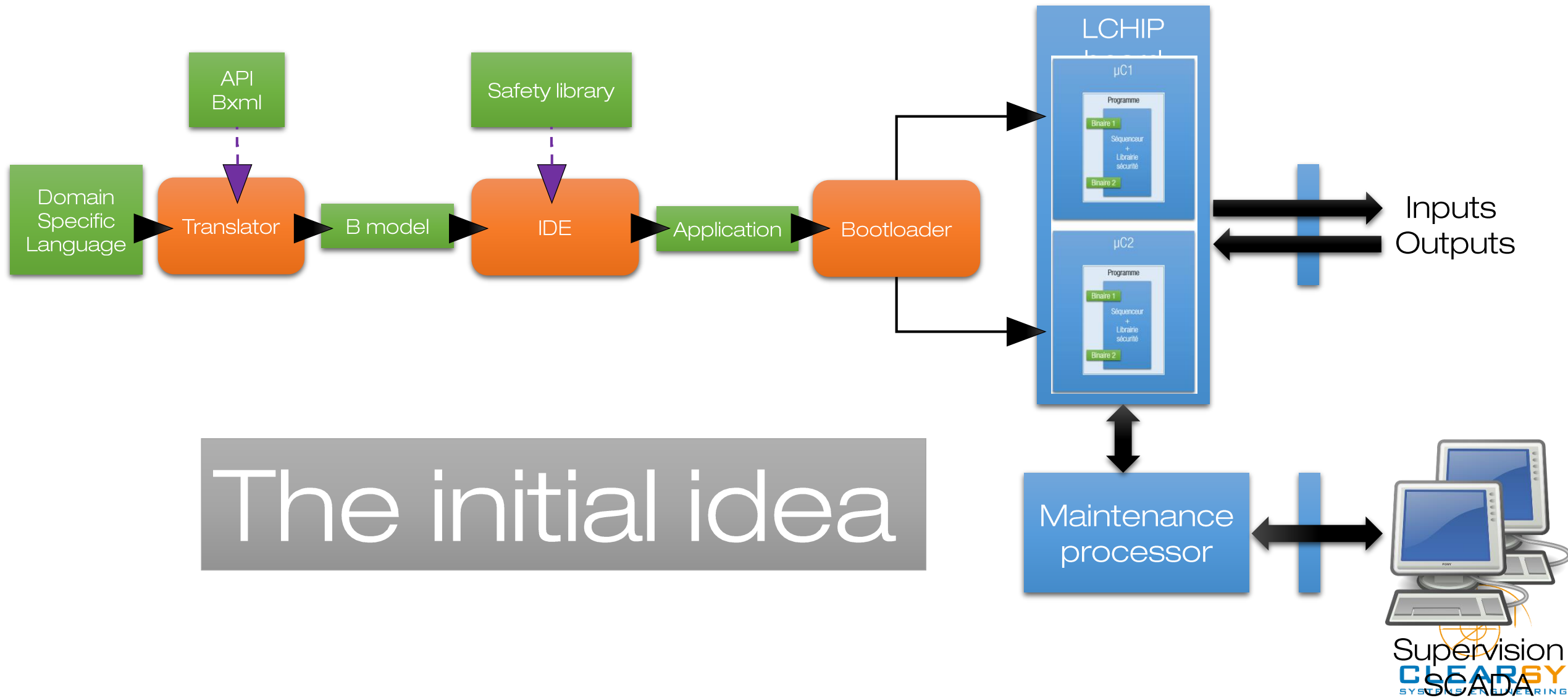
To beam up a function on 2 microcontrolers



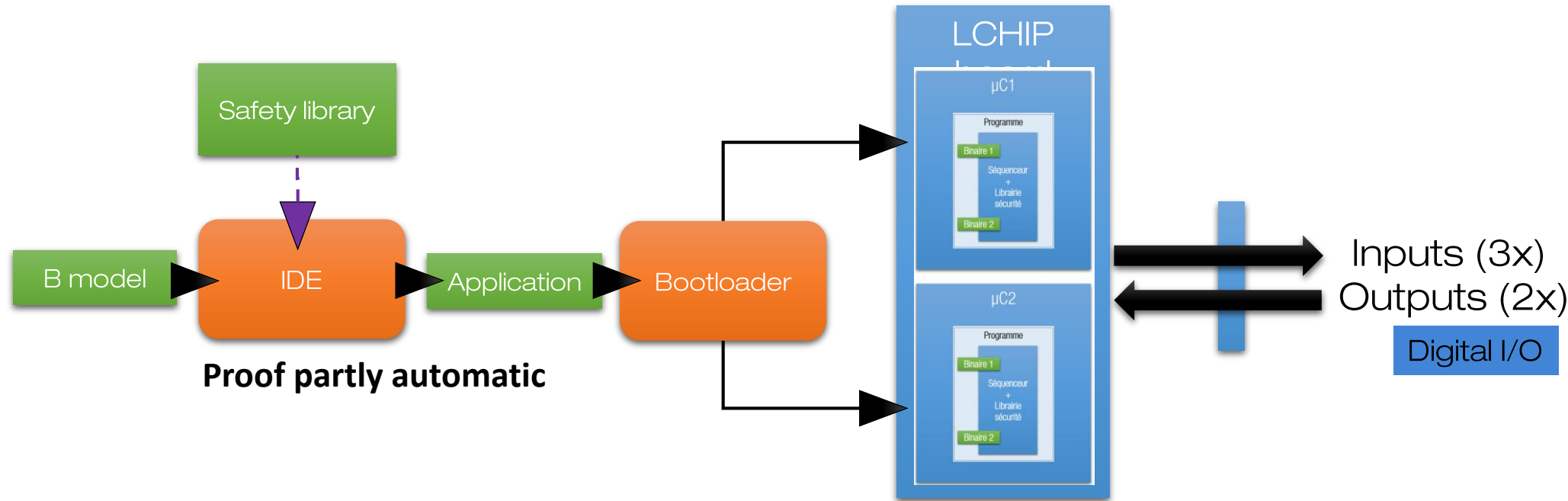
The Software Factory



Roadmap

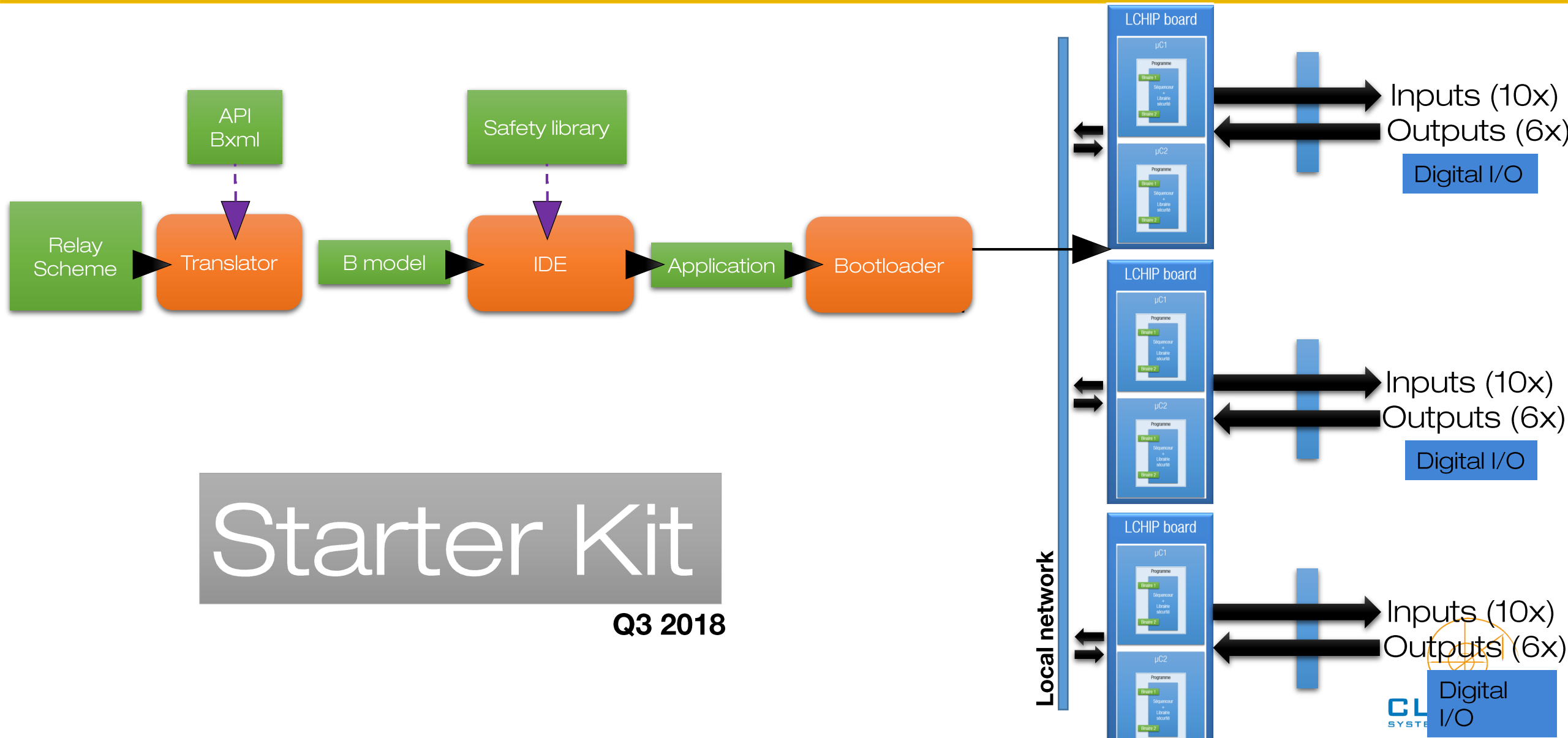


Roadmap

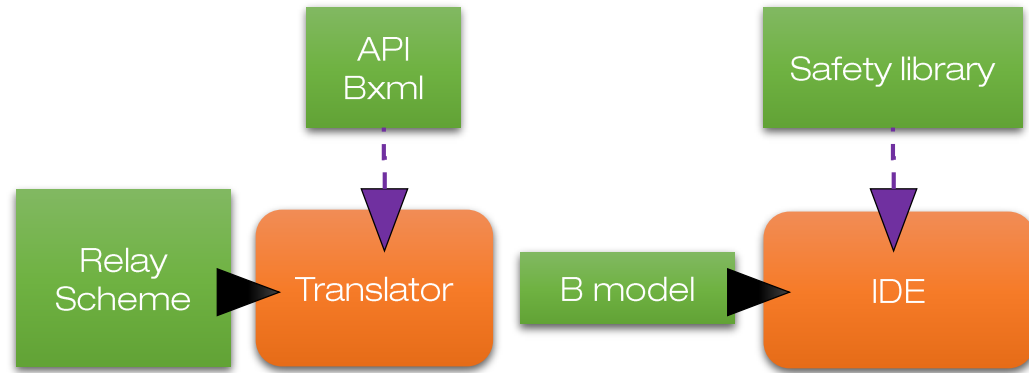


Starter Kit SK₀

Roadmap



Roadmap

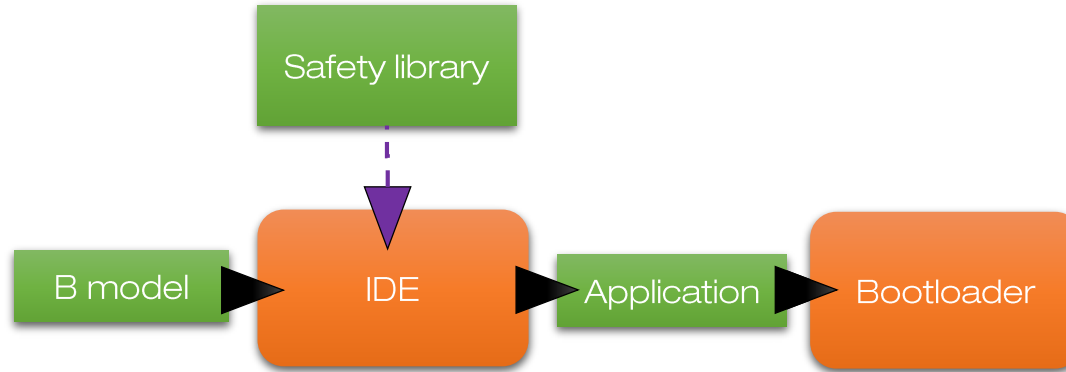


Translation from relay scheme to B

Feedback from the IDE to the initial relay scheme model

- Unsupported B construct
- Unproved proof obligation
- Any information to help the engineer to fine tune the DSL model

Roadmap

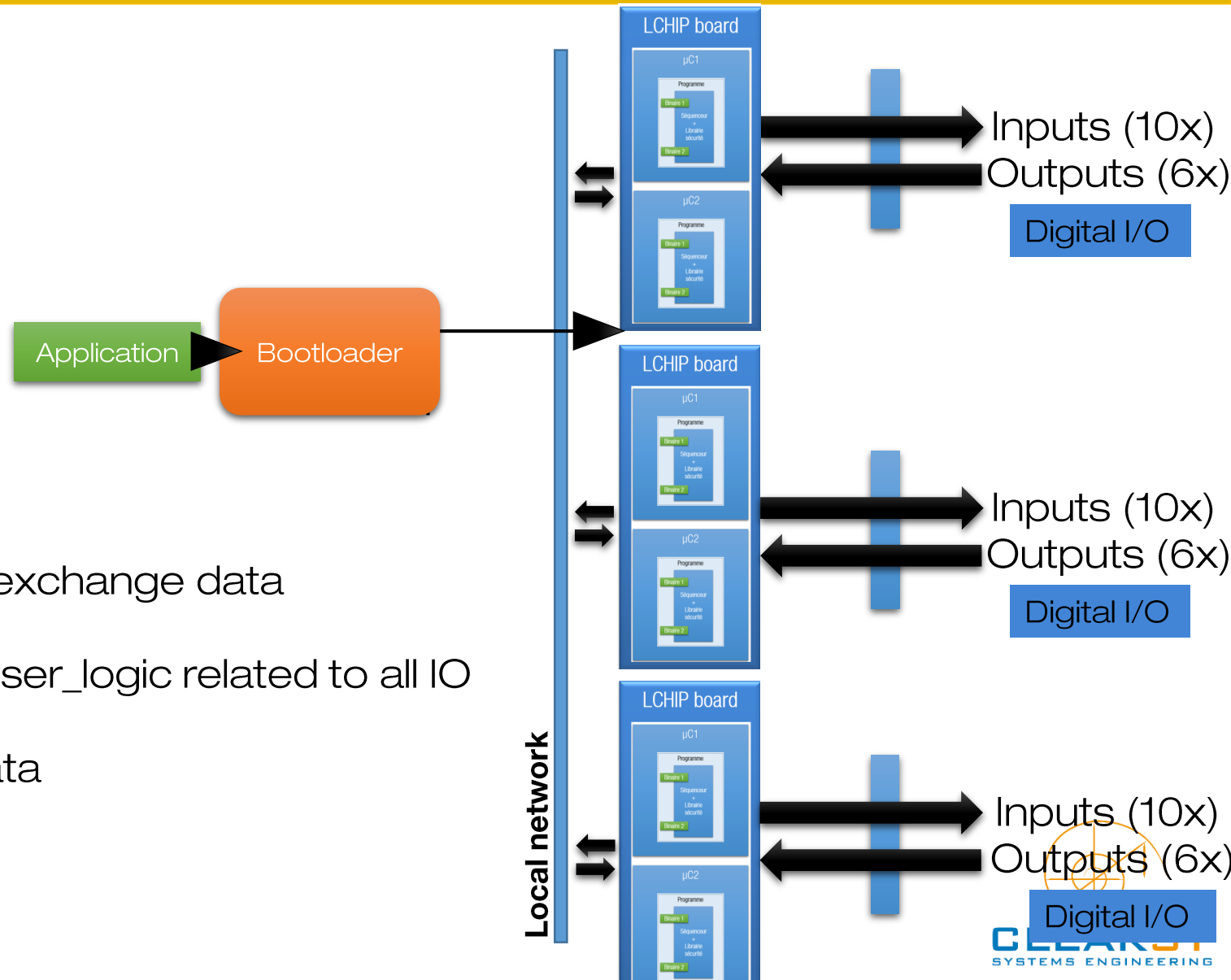


Improved, more resilient IDE

Better integration with the application generation toolchain

Improved proof performances

Roadmap



More I/O per board

Boards to connect on a local network to exchange data

Every LCHIP board compute the overall user_logic related to all IO

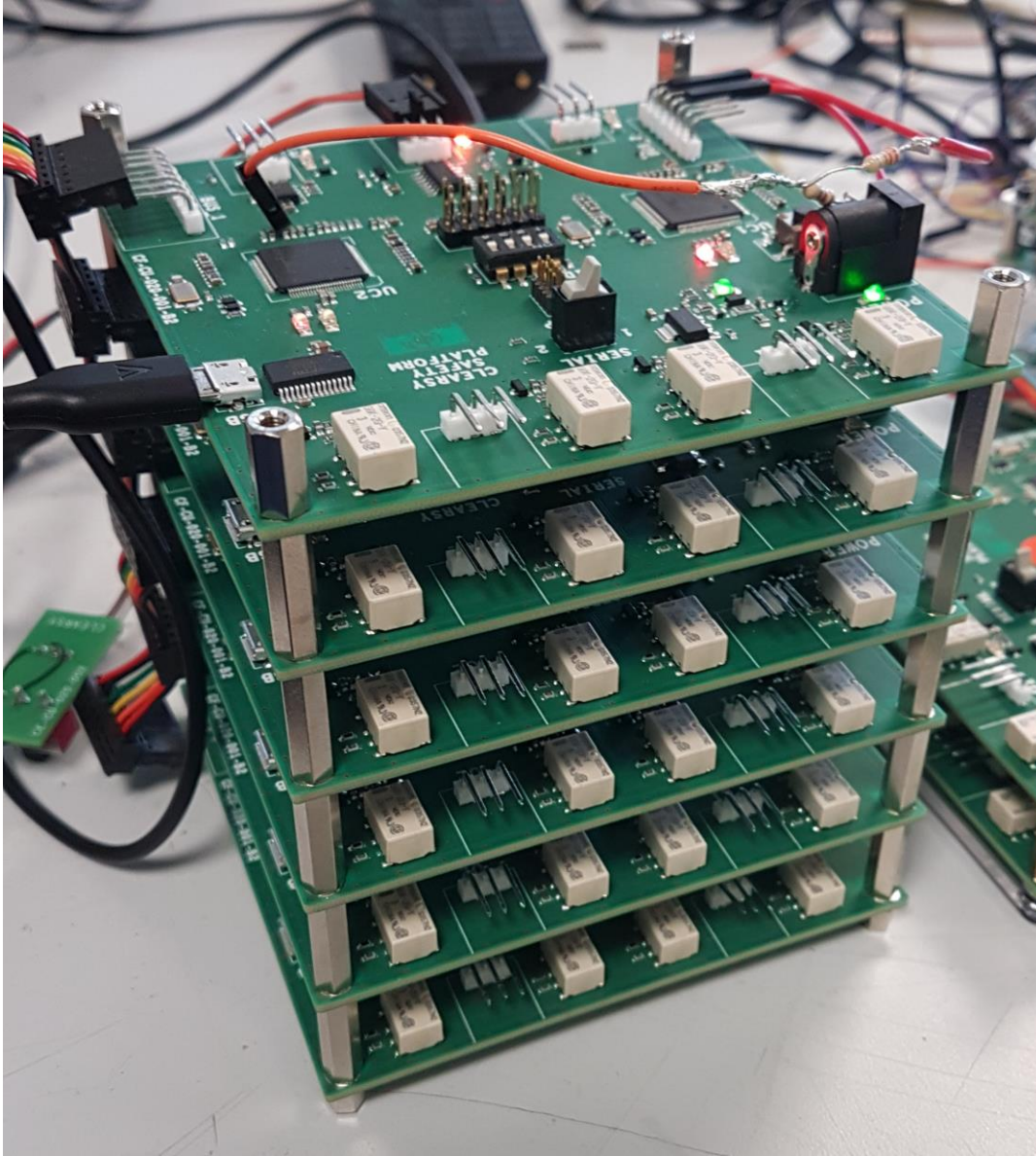
B model generated from configuration data

Roadmap



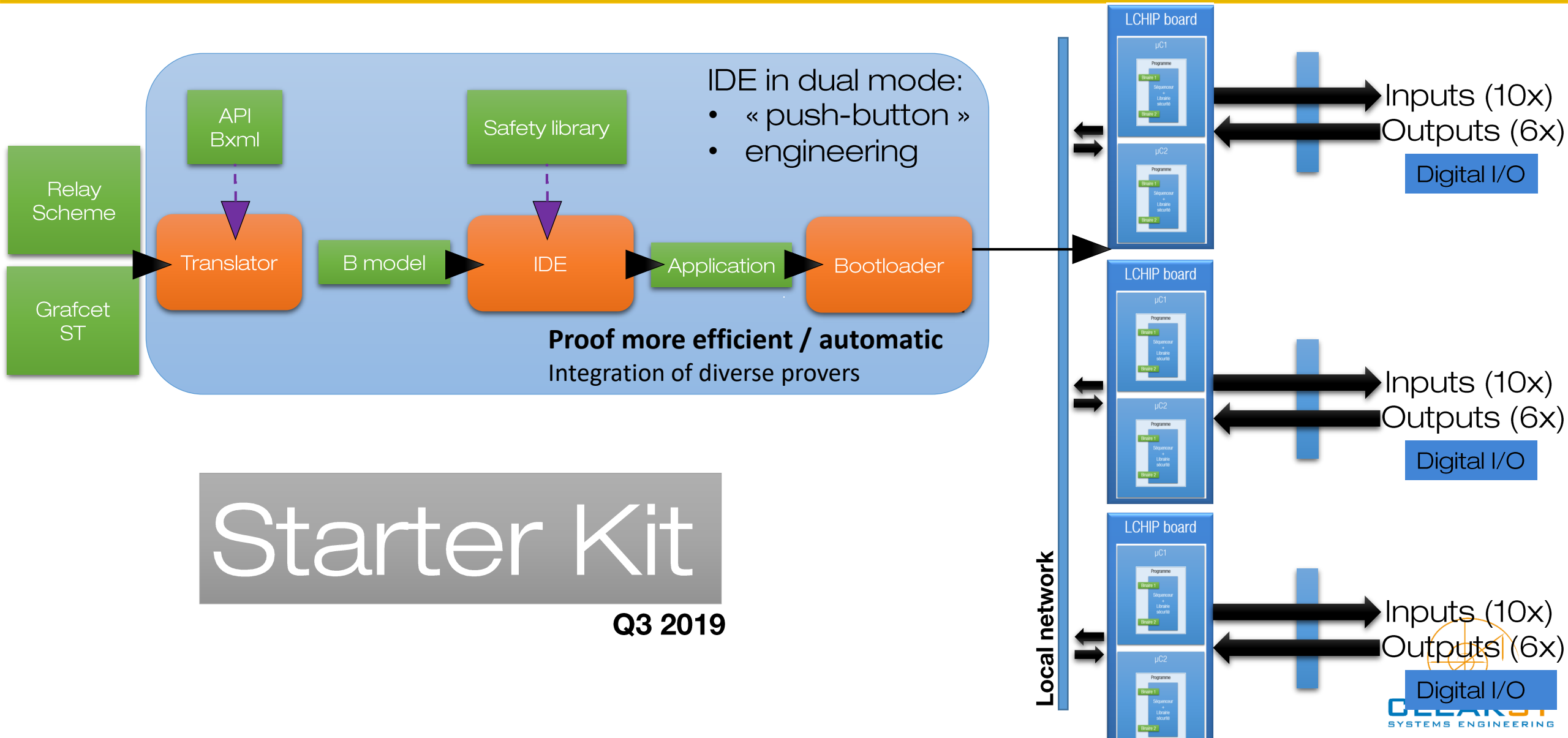
2 boards connected
sharing inputs status

Roadmap



More boards ...

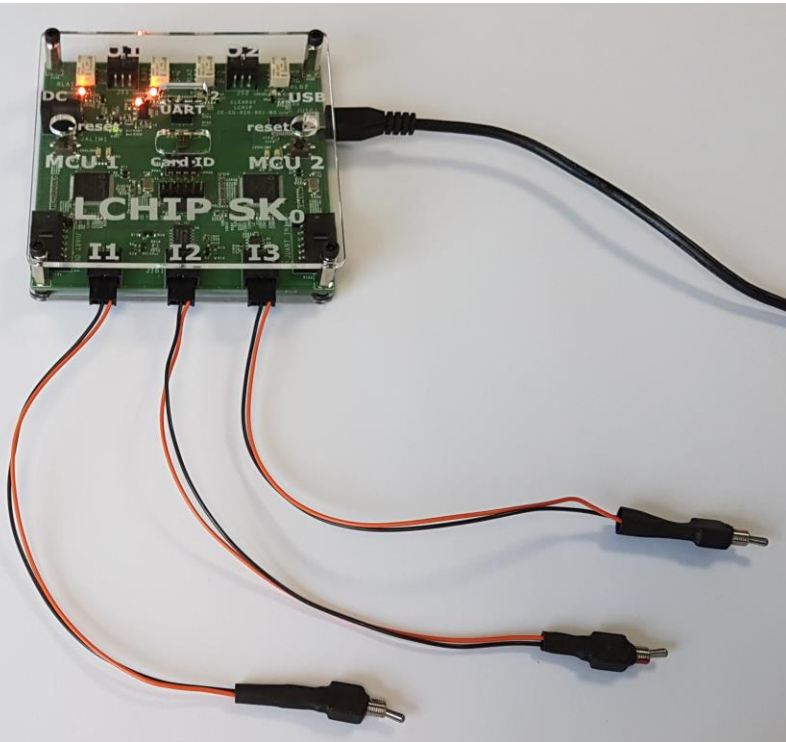
Roadmap



Starter Kit

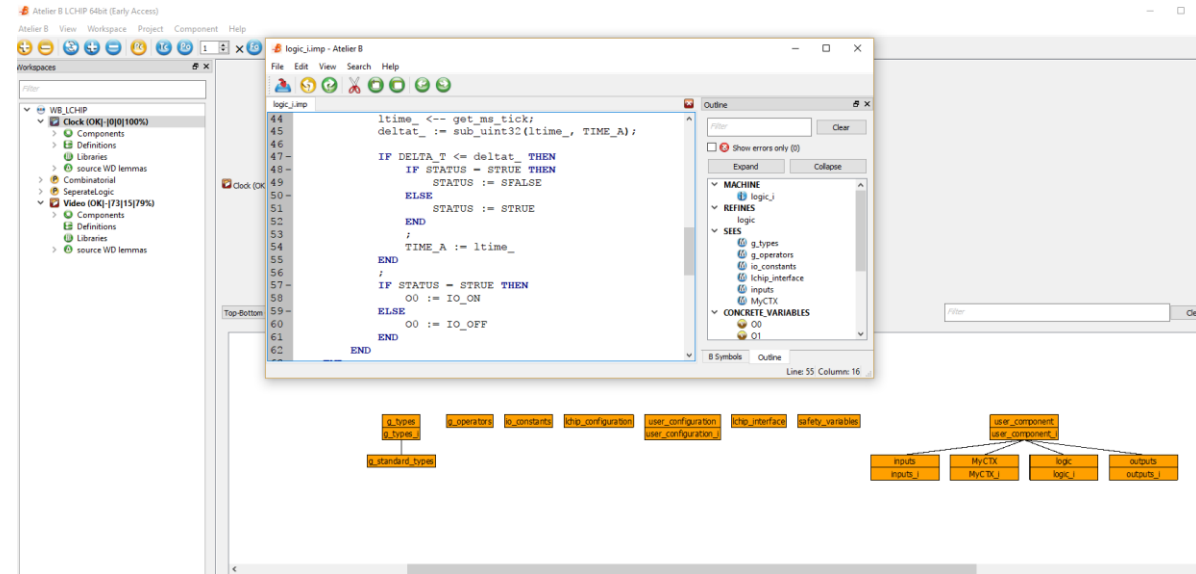
Q3 2019

Starter Kit SK₀



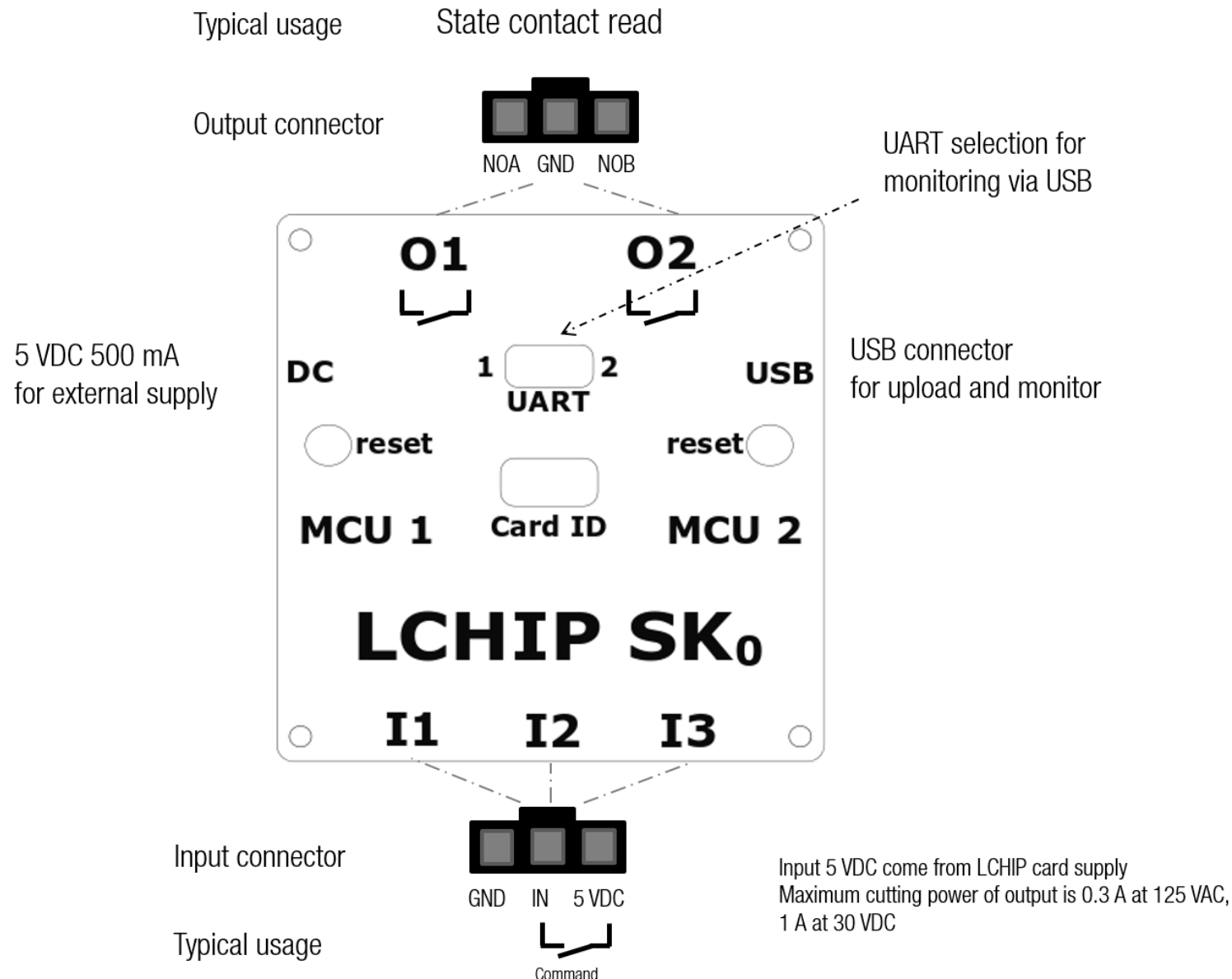
USB-C connector
Power, upload, monitor

Switches
Simulate digital inputs



Atelier B 4.4 LCHIP Early access
Model, prove, compile, upload

Starter Kit SK₀



A complete transaction

DEMO

A complete transaction

- Installing the IDE

Copy the the content of the USB in a directory that contains no space nor no special character. Execute Register LCHIP.cmd
- Starting the IDE

Execute startAB.cmd. For the first execution, change Atelier B/preferences/langage to english/us. Quit then restart to get the english GUI
- Creating a project

Click the yellow + button, select software development, give a name, press next then finish. Double-click on the project name to open it.
- Importing the LCHIP infrastructure

Select an open project. Select Import LCHIP API. Press Fermer. Press F5 to update the project status to see added components.
- Contributing

Double-click on a component. Modify it. Ctrl+S to save. Select the component then press the blue TC button to operate the typecheck.
- Proving

Select the component then press the blue F0 button to operate the proof.
- Compiling

Select the LCHIP project then select compile LCHIP.
- Uploading

Select upload. Select connect on the uploader window. Select Erase-Program-Verify. Trigger one reset button on the board. When the device is ready, trigger one reset button to leave bootload mode.
- Monitoring

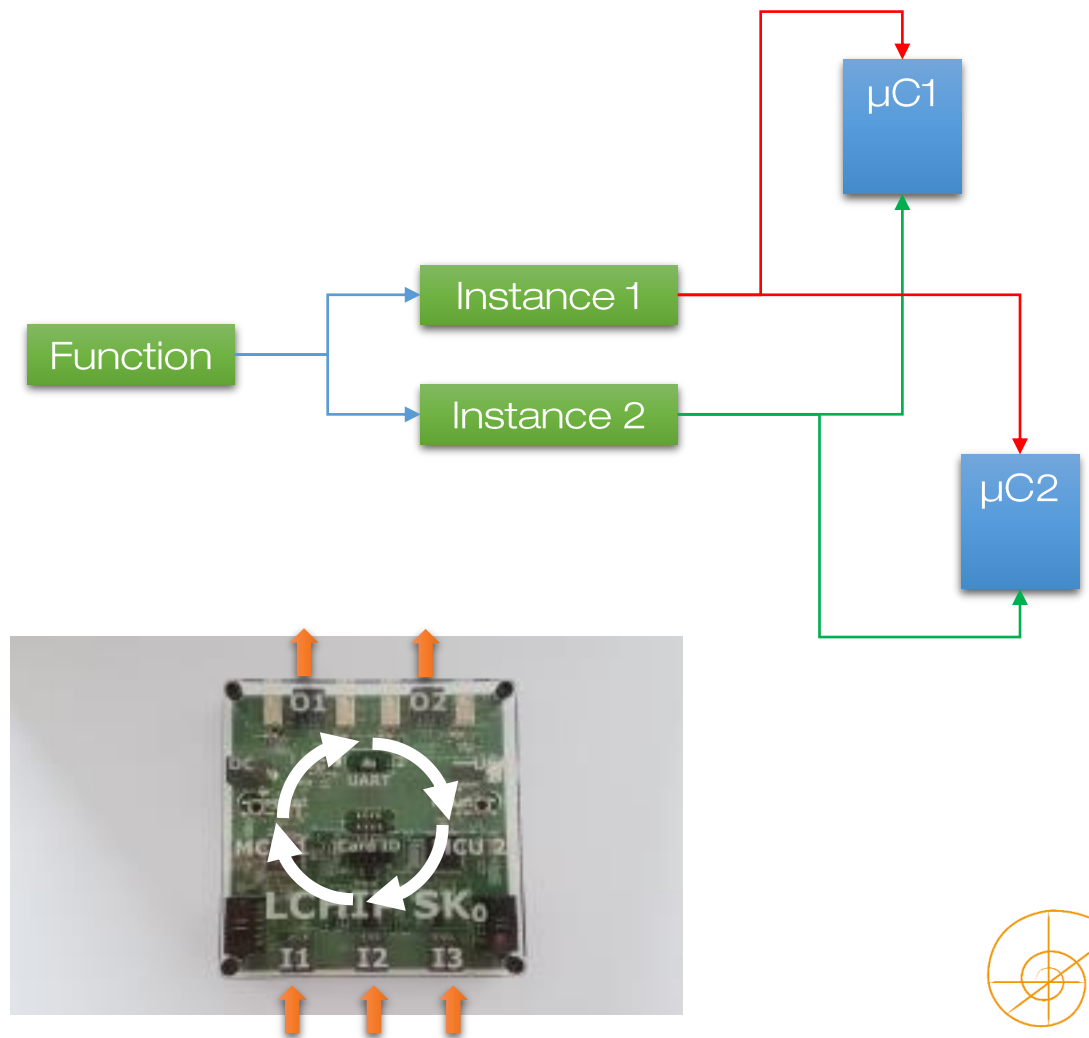
Select monitor on the uploader window. Disconnect to stop monitoring

Programming

The programming model

The developer is focused on developing a function

- Independently of its transformation and distributed execution
- Can read inputs
- Can perform computations using a subset of the B language
- Can modify the outputs
- Can read the current time since the CU started



The programming model

The execution is cyclic

- The function is executed regularly as often as possible similar to arduino programming (setup(), loop())
- No underlying operating system
- No interrupt
- No predefined cycle time (if outputs are not set and cross read every 50 ms, SK₀ enters panic mode)
- No delay()
- Inputs are values captured at the beginning of a cycle (instantaneous values)
- Outputs are maintained from one cycle to another

```
init();  
  
while (1) {  
    instance1();  
    instance2();  
}
```

The programming model

LCHIP takes care of all safety verifications

If a divergent behaviour is detected

- one of the 4 software instances behaves differently
- one UC behaves differently

or if a structural error is detected

- bad CRC on memory
- UC unable to execute an operation properly
- etc.

then the detecting UC reboots

Verification intra-MCU
13 000 per second

Verification inter-MCU
48 per second

measured on Clock example

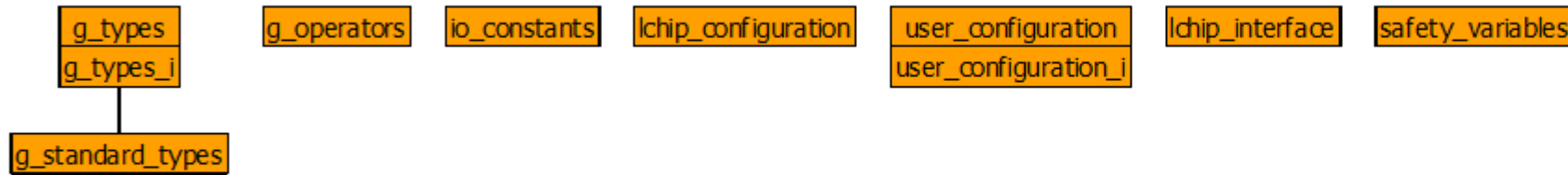
Verification deferred
over several cycles

A LCHIP Project

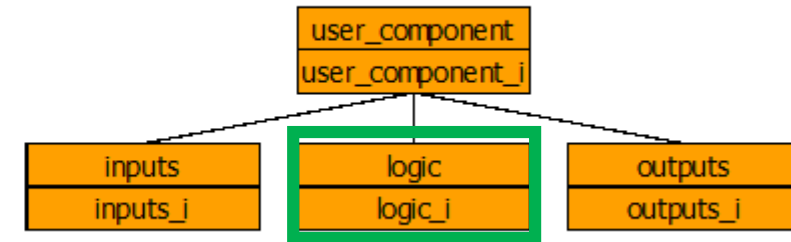
is a B project

- is generated automatically from board configuration (# IOs, naming)
- in SK_0 , more information than required is exposed

It contains



The interface with the safety library



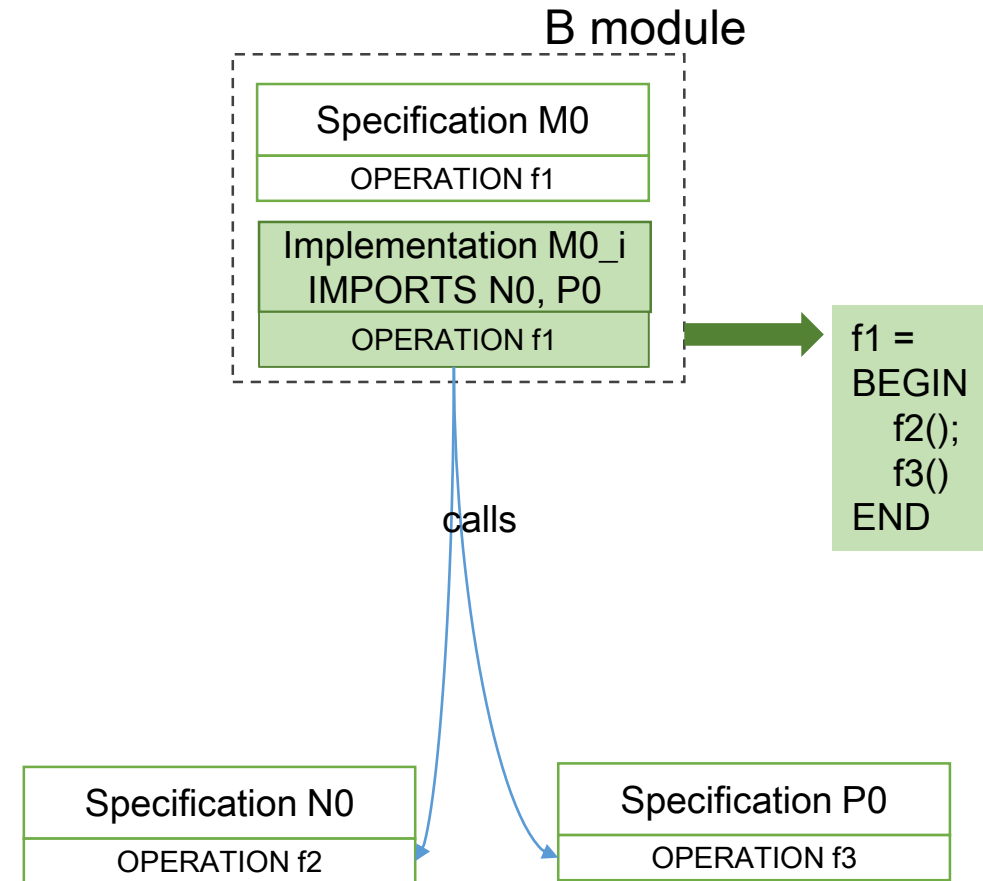
Components to modify

The model of the function



Reminder: Models Architecture

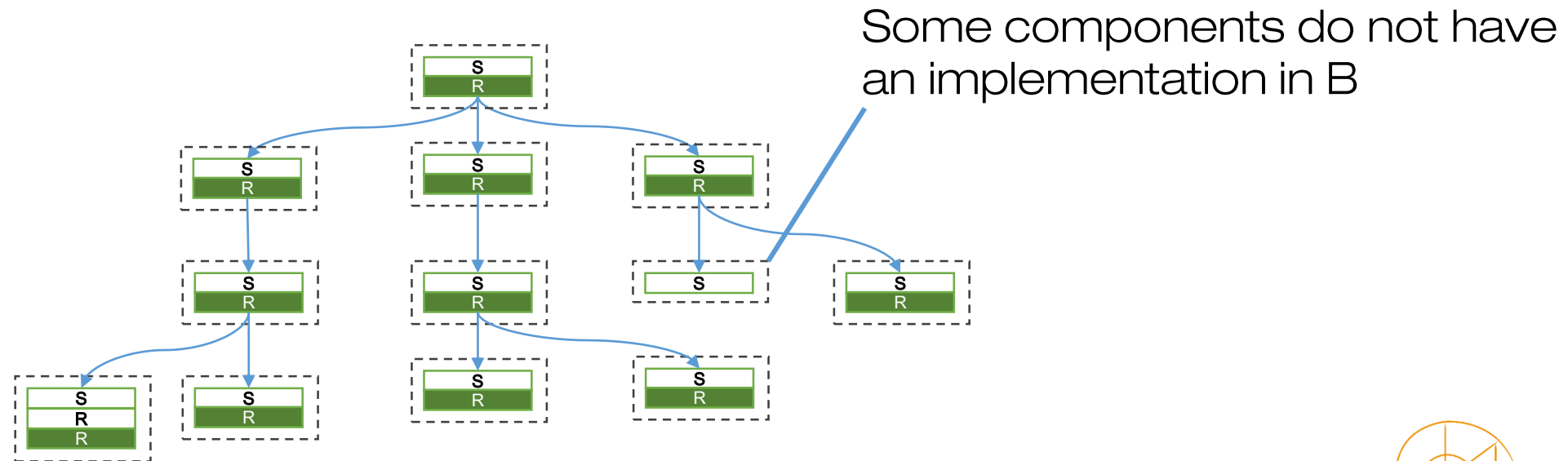
- One operation cannot call other operations from the same implementation
- One operation can call operations defined in other machines
- These machines have to be imported in an implementation
- Variables defined in imported machines have to be different: a variable cannot be modified in 2 components



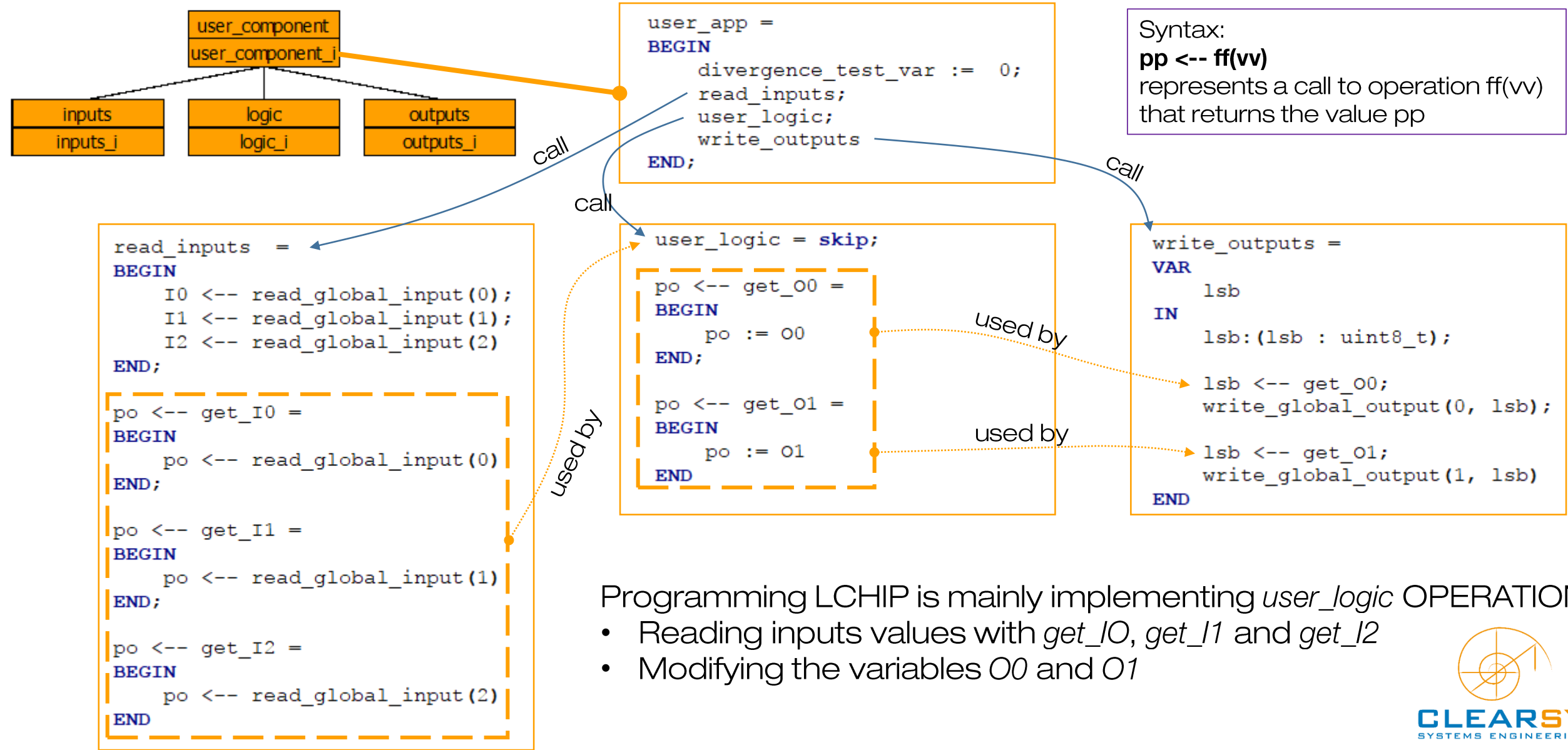
Reminder: Models Architecture

IMPORTS have to be applied iteratively to obtain the target decomposition

The decomposition graph should be a tree



Function model



Introduction to B: variables declaration

specification

ABSTRACT_VARIABLES

```
O0,  
O1
```

: means
« belongs to »

INVARIANT

```
O0 : uint8_t &  
O1 : uint8_t
```

|| means « in parallel », « at the
same time »

INITIALISATION

```
O0 :: uint8_t ||  
O1 :: uint8_t
```

:: means « any value within »

implementation

```
// pragma SAFETY_VARS
```

Mandatory

Contains variables
that will be verified

CONCRETE_VARIABLES

```
O0,  
O1,  
TIME_A,  
STATUS
```

} Variables local to
implementation

INVARIANT

```
O0 : uint8_t &  
O1 : uint8_t &  
TIME_A : uint32_t &  
STATUS : uint8_t
```

INITIALISATION

```
O0 := IO_OFF;  
O1 := IO_OFF;  
TIME_A := 0;  
STATUS := SFALSE
```



Introduction to B: constants declaration

specification

```
CONCRETE_CONSTANTS  
  DELTA_T
```

```
PROPERTIES  
  DELTA_T : uint32_t
```

implementation

```
// pragma CONSTANTS — Mandatory Contains constants  
that will be verified
```

Important

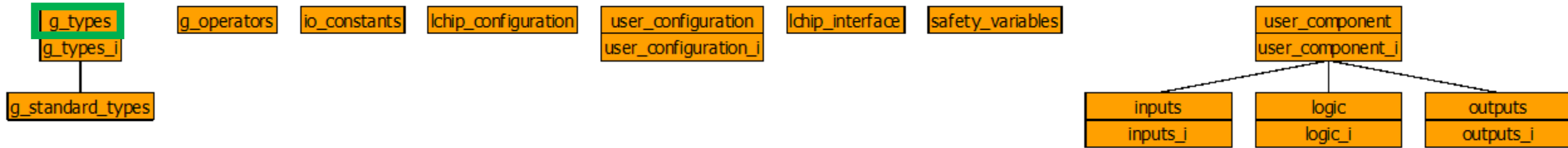
A model cannot contain both variables
and constants

```
VALUES  
  DELTA_T = 1000 // 1000 ms == 1s
```

Value that should enforce the
properties



Introduction to B: types



CONCRETE CONSTANTS

```
uint32_t ,
uint16_t ,
uint8_t ,
```

Everything is either 8, 16 or 32 bits

```
STRUE ,
SFALSE ,
MAX_UINT32 ,
MAX_UINT16 ,
MAX_UINT8
```

Boolean values TRUE and FALSE coded on 8 bits

The real values for STRUE and SFALSE are not displayed but we know that

PROPERTIES

```
uint32_t = 0 .. 4294967295 &
uint16_t = 0 .. 65535 &
uint8_t = 0 .. 255 &
```

```
MAX_UINT32 : uint32_t &
MAX_UINT16 : uint16_t &
MAX_UINT8 : uint8_t &
```

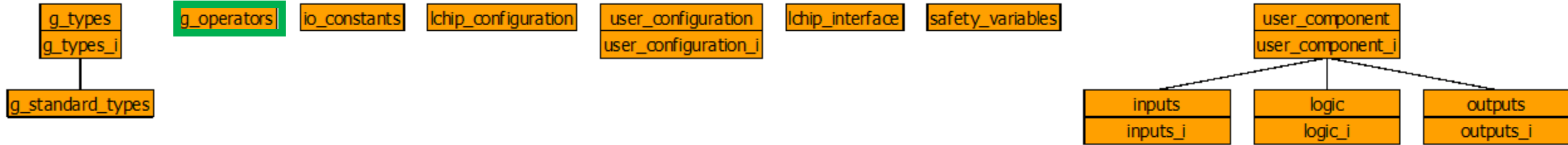
```
STRUE : uint8_t &
SFALSE : uint8_t &
```

```
MAX_UINT32 = 4294967295 &
MAX_UINT16 = 65535 &
MAX_UINT8 = 255 &
```

```
STRUE : 0 .. MAX_UINT8 &
SFALSE : 0 .. MAX_UINT8 &
```

```
STRUE /= SFALSE &
SBOOL = { STRUE , SFALSE } &
STRUE <= 2 &
SFALSE <= 2 &
```

Introduction to B: unsigned int operators



CONCRETE CONSTANTS

```
bitwise_not_uint32,
bitwise_and_uint32,
bitwise_xor_uint32,
bitwise_not_uint16,
bitwise_and_uint16,
bitwise_xor_uint16,
bitwise_or_uint16,
bitwise_not_uint8,
bitwise_and_uint8,
bitwise_xor_uint8,
bitwise_or_uint8,
```

```
add_uint32,
sub_uint32,
mul_uint32,
add_uint16,
sub_uint16,
mul_uint16,
add_uint8,
sub_uint8,
mul_uint8
```

Builtin operators

PROPERTIES

```
bitwise_not_uint32 : uint32_t --> uint32_t &
bitwise_and_uint32 : uint32_t * uint32_t --> uint32_t &
bitwise_xor_uint32 : uint32_t * uint32_t --> uint32_t &
bitwise_not_uint16 : uint16_t --> uint16_t &
bitwise_and_uint16 : uint16_t * uint16_t --> uint16_t &
bitwise_xor_uint16 : uint16_t * uint16_t --> uint16_t &
bitwise_or_uint16  : uint16_t * uint16_t --> uint16_t &
bitwise_not_uint8  : uint8_t  --> uint8_t  &
bitwise_and_uint8  : uint8_t  * uint8_t  --> uint8_t  &
bitwise_xor_uint8  : uint8_t  * uint8_t  --> uint8_t  &
bitwise_or_uint8   : uint8_t  * uint8_t  --> uint8_t  &
```

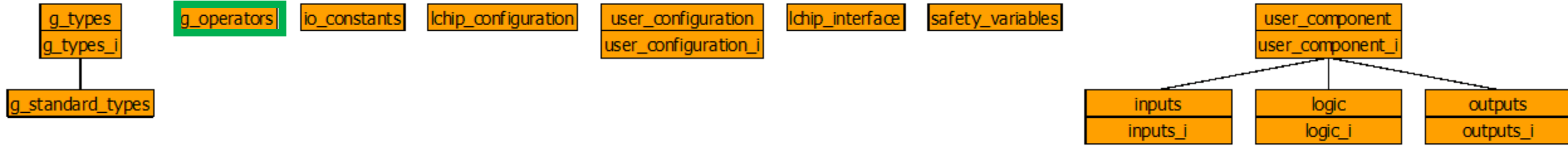
`bitwise_not_uint32 : uint32_t --> uint32_t`

total function that associates a 32-bit unsigned integer to any 32-bit unsigned integer

`bitwise_and_uint32 : uint32_t * uint32_t --> uint32_t`

total function with two 32-bit unsigned integer parameters

Introduction to B: unsigned int operators



CONCRETE CONSTANTS

```
bitwise_not_uint32,  
bitwise_and_uint32,  
bitwise_xor_uint32,  
bitwise_not_uint16,  
bitwise_and_uint16,  
bitwise_xor_uint16,  
bitwise_or_uint16,  
bitwise_not_uint8,  
bitwise_and_uint8,  
bitwise_xor_uint8,  
bitwise_or_uint8,
```

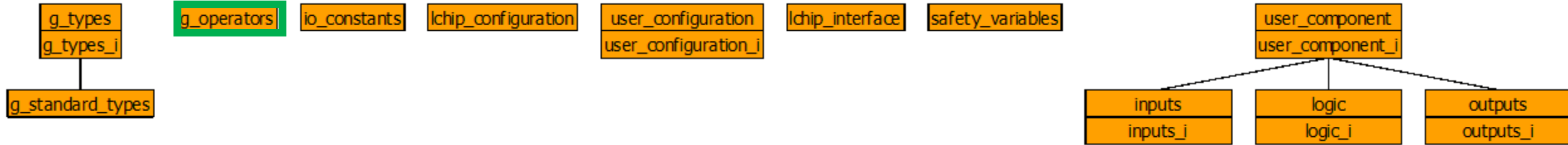
```
add_uint32,  
sub_uint32,  
mul_uint32,  
add_uint16,  
sub_uint16,  
mul_uint16,  
add_uint8,  
sub_uint8,  
mul_uint8
```

Builtin operators

```
add_uint32 : uint32_t * uint32_t --> uint32_t &  
sub_uint32 : uint32_t * uint32_t --> uint32_t &  
mul_uint32 : uint32_t * uint32_t --> uint32_t &  
add_uint16 : uint16_t * uint16_t --> uint16_t &  
sub_uint16 : uint16_t * uint16_t --> uint16_t &  
mul_uint16 : uint16_t * uint16_t --> uint16_t &  
add_uint8 : uint8_t * uint8_t --> uint8_t &  
sub_uint8 : uint8_t * uint8_t --> uint8_t &  
mul_uint8 : uint8_t * uint8_t --> uint8_t &
```



Introduction to B: unsigned int operators



```
add_uint32 = % (x1, x2) . (x1 : uint32_t & x2 : uint32_t | (x1 + x2) mod (MAX_UINT32 + 1)) &
sub_uint32 = % (x1, x2) . (x1 : uint32_t & x2 : uint32_t | (x1 - x2 + MAX_UINT32 + 1) mod (MAX_UINT32 + 1)) &
mul_uint32 = % (x1, x2) . (x1 : uint32_t & x2 : uint32_t | (x1 * x2) mod (MAX_UINT32 + 1)) &
add_uint16 = % (y1, y2) . (y1 : uint16_t & y2 : uint16_t | (y1 + y2) mod (MAX_UINT16 + 1)) &
sub_uint16 = % (y1, y2) . (y1 : uint16_t & y2 : uint16_t | (y1 - y2 + MAX_UINT16 + 1) mod (MAX_UINT16 + 1)) &
mul_uint16 = % (y1, y2) . (y1 : uint16_t & y2 : uint16_t | (y1 * y2) mod (MAX_UINT16 + 1)) &
add_uint8 = % (y1, y2) . (y1 : uint8_t & y2 : uint8_t | (y1 + y2) mod (MAX_UINT8 + 1)) &
sub_uint8 = % (y1, y2) . (y1 : uint8_t & y2 : uint8_t | (y1 - y2 + MAX_UINT8 + 1) mod (MAX_UINT8 + 1)) &
mul_uint8 = % (y1, y2) . (y1 : uint8_t & y2 : uint8_t | (y1 * y2) mod (MAX_UINT8 + 1)) &
```

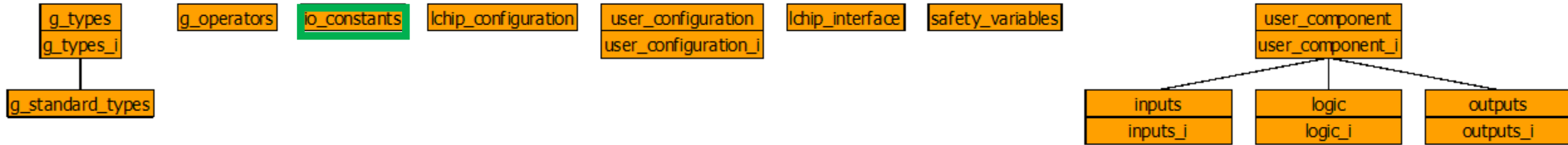
```
add_uint32 = % (x1, x2) . (x1 : uint32_t & x2 : uint32_t | (x1 + x2) mod (MAX_UINT32 + 1))
```

is a λ function

that takes two 32-bit
unsigned integer parameters

and returns the sum of the values
modulo $\text{MAX_UINT32} + 1$

Introduction to B: inputs/outputs



ABSTRACT CONSTANTS

```
TIME,  
IO_STATE
```

inputs and outputs state

CONCRETE CONSTANTS

```
IO_ON,  
IO_OFF
```

values used by digital inputs and outputs

PROPERTIES

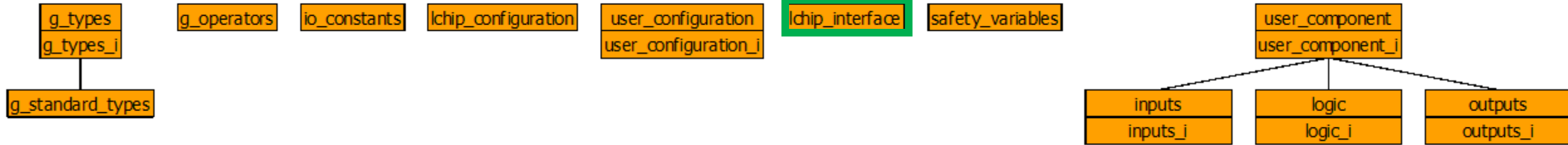
```
TIME = uint32_t &  
IO_STATE = uint8_t &  
  
IO_ON : uint8_t &  
IO_OFF : uint8_t &  
IO_ON /= IO_OFF &  
IO_ON : IO_STATE &  
IO_OFF : IO_STATE
```

coded on 8 bits

Verification

If a digital output is valued with a value different from IO_ON or IO_OFF then SK₀ stops in error mode

Introduction to B: inputs/outputs



```
out <-- get_ms_tick =
PRE
  out : uint32_t
THEN
  out := ms_tick
END
```

————— returns the number of milliseconds since the last reset

Important

SK₀ resets when the ms_tick reaches its upper bound
i.e. every 49.7 days

Introduction to B: operations

Operations are populated with substitutions

Available substitutions in specification are different from the ones available in implementation

specification

Express the properties that the variables comply with when the operation is completed independently from the algorithm implemented (*post-condition*)

To simplify, always use « becomes such that substitutions »

```
user_logic =  
  BEGIN  
    O0, O1 : (  
      O0 : uint8_t & } Typing (mandatory)  
      O1 : uint8_t & } Constraints (optional)  
      not(O0 = O1)  
    )  
  END;
```

Introduction to B: operations

implementation

`user_logic = skip;` — do nothing

```
user_logic =  
BEGIN  
  O0 := IO_ON;  
  O1 := IO_OFF;  
END;
```

— valuations in sequence

```
user_logic =  
BEGIN  
  IF Var8 = 0 THEN  
    O0 := IO_ON  
  ELSE  
    O1 := IO_ON  
  END  
END;
```

— IF THEN ELSE

Important

Only single condition (no conjunction nor disjunction)
= < <= operators only

```
user_logic =  
BEGIN  
  VAR time_ IN  
    time_ : (time_ : uint32_t);  
    time_ <-- get_ms_tick;  
  IF 2000 <= time_ THEN  
    O1 := IO_ON  
  END  
END  
END;
```

— Local variables declaration
— Operation call

Important

Local variables have to be typed first using
« becomes such that » substitution

Introduction to B: user_logic

specification

```
user_logic = skip;
```

skip means « do not alter the variables of the model »

```
MACHINE
  logic

SEES
  g_types,
  g_operators,
  io_constants,
  lchip_interface

ABSTRACT_VARIABLES
  O1,
  O2

INVARIANT
  O1 : uint8_t &
  O2 : uint8_t

INITIALISATION
  O1 :: uint8_t ||
  O2 :: uint8_t

OPERATIONS
  user_logic = skip;

  po <-- get_O1 =
  PRE
    po : uint8_t
  THEN
    po := O1
  END;

  po <-- get_O2 =
  PRE
    po : uint8_t
  THEN
    po := O2
  END
END
```

implementation

```
user_logic = skip;
```

Minimum example:

- do nothing; outputs remain in their initial state (INITIALISATION)

```
IMPLEMENTATION logic_i
REFINES logic

SEES
  g_types,
  g_operators,
  io_constants,
  lchip_interface,
  inputs

  // pragma SAFETY_VARS

CONCRETE_VARIABLES
  O1,
  O2

INVARIANT
  O1 : uint8_t &
  O2 : uint8_t

INITIALISATION
  O1 := IO_OFF;
  O2 := IO_OFF

OPERATIONS
  user_logic = skip;

  po <-- get_O1 =
  BEGIN
    po := O1
  END;

  po <-- get_O2 =
  BEGIN
    po := O2
  END
END
```

Introduction to B: user_logic

specification

```
user_logic =  
BEGIN  
  O0 :: uint8_t ||  
  O1 :: uint8_t  
END
```

— O0 and O1 belong to their type

```
user_logic =  
BEGIN  
  O0, O1 : (  
    O0 : uint8_t &  
    O1 : uint8_t &  
    not(O0 = O1)  
  )  
END
```

:() means « becomes such that »

— O0 and O1 belong to their type and O0 is different from O1

```
user_logic =  
BEGIN  
  O0 := IO_ON ||  
  O1 := IO_OFF  
END
```

— Set O0 and reset O1

implementation

```
user_logic =  
BEGIN  
  O0 := IO_ON;  
  O1 := IO_OFF  
END
```

— Set O0 then reset O1

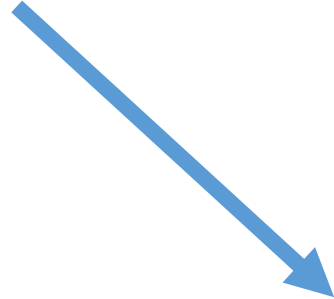
« then » is related to the valuation of O0 regarding O1
O0 and O1 will be positioned at the same time at the end of the cycle

Example #1: combinatorial hello world

Example #1: combinatorial

$O_0 = I_0 \text{ and } I_1 \text{ and } I_2$

$O_1 = \text{not}(O_0)$



```
user_logic =  
BEGIN
```

```
    VAR i0_, i1_, i2_ IN
```

```
        i0_ : (i0_ : uint8_t);
```

```
        i1_ : (i1_ : uint8_t);
```

```
        i2_ : (i2_ : uint8_t);
```

} Local variables are typed first

```
        i0_ <-- get_I0;
```

```
        i1_ <-- get_I1;
```

```
        i2_ <-- get_I2;
```

} Local variables are valued

```
        O0 <-- triAND(i0_, i1_, i2_); /* O0 is ON iff I0, I1 & I2 are ON */
```

```
        O1 <-- negIO(O0) /* O1 is the opposite of O0 */
```

```
    END
```

```
END
```

```
;
```

Variables are valued with LOCAL_OPERATIONS

Example #1: combinatorial

LOCAL_OPERATIONS

```
res <-- triAND(v1, v2, v3) =  
PRE  
  v1: uint8_t & v2: uint8_t & v3: uint8_t  
THEN  
  res :: uint8_t  
END  
;  
res <-- negIO(val) =  
PRE  
  val : uint8_t  
THEN  
  res :: uint8_t  
END
```

} Input parameters have to be type first in the precondition clause

Syntax: **PRE** predicates **THEN** substitution **END**

Operations specified in LOCAL_OPERATIONS have to be implemented in OPERATIONS

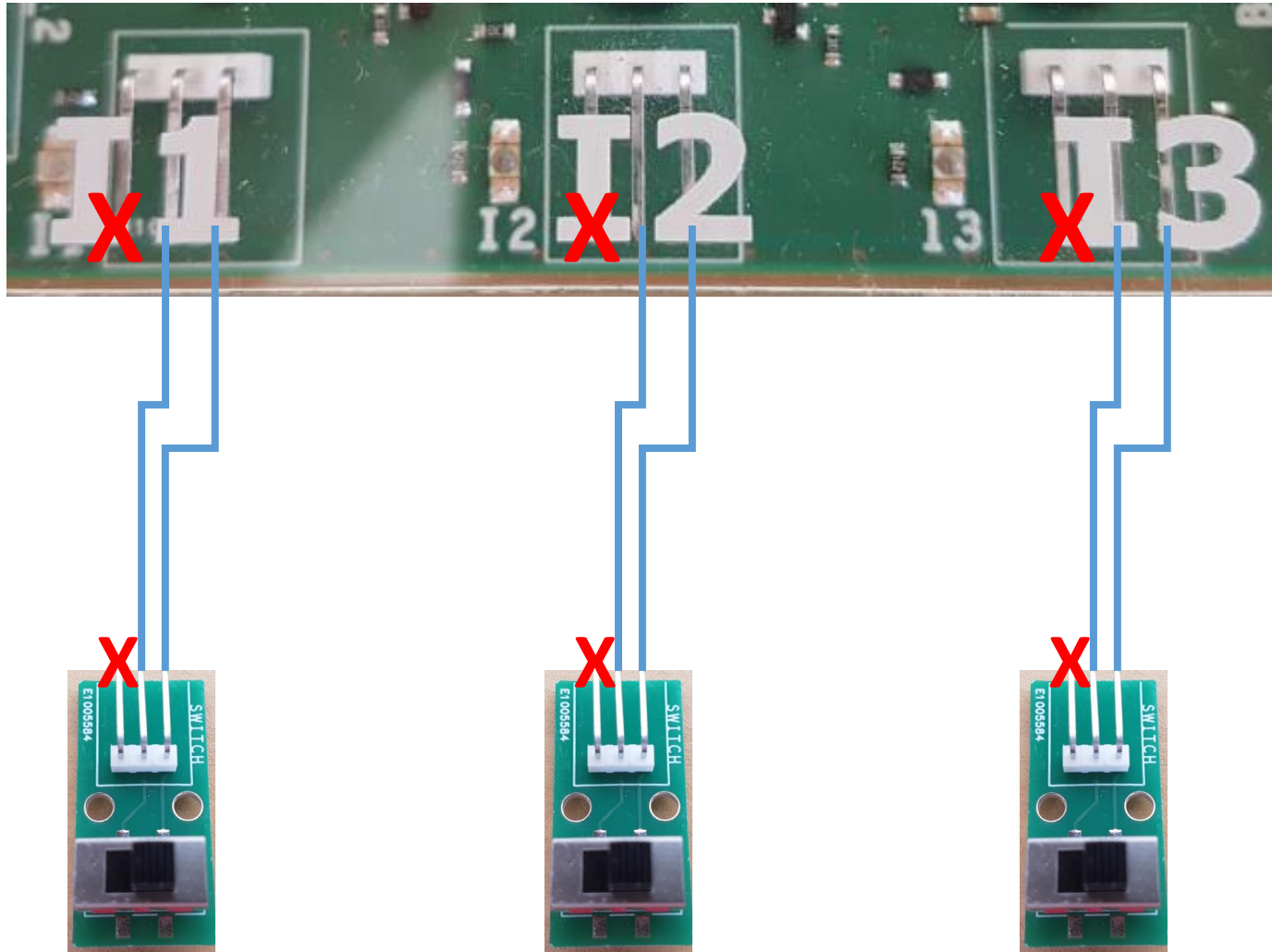
OPERATIONS

```
res <-- triAND(v1, v2, v3) = /* AND over 3 values */  
BEGIN  
  res :( res : uint8_t);  
  res := IO_OFF;  
  IF v1 = IO_ON THEN  
    IF v2 = IO_ON THEN  
      IF v3 = IO_ON THEN  
        res := IO_ON  
      END  
    END  
  END  
END  
END  
END
```

} Output parameters have to be type first



Example #1: combinatorial



How to connect switches to the board

Example #1: combinatorial

Your turn:

- **Instead of a AND, program a OR over the 3 inputs**
- **Model:**
 - Rename triAND operation in triOR
 - Modify triOR and user_logic implementations
- **Prove:**
 - Ctrl+A (component view) to select all components, press F0 to start type check, PO generation and proof in sequence
- **Compile:**
 - Right click on the project (left pane) then select « Compile LCHIP M »
- **Upload:**
 - Connect your board, click on « upload », click on « connect », click on « erase program verify », reset your board, wait for « device ready », reset your board
- **Check** with your buttons that the function is correctly implemented

Example #2: clock

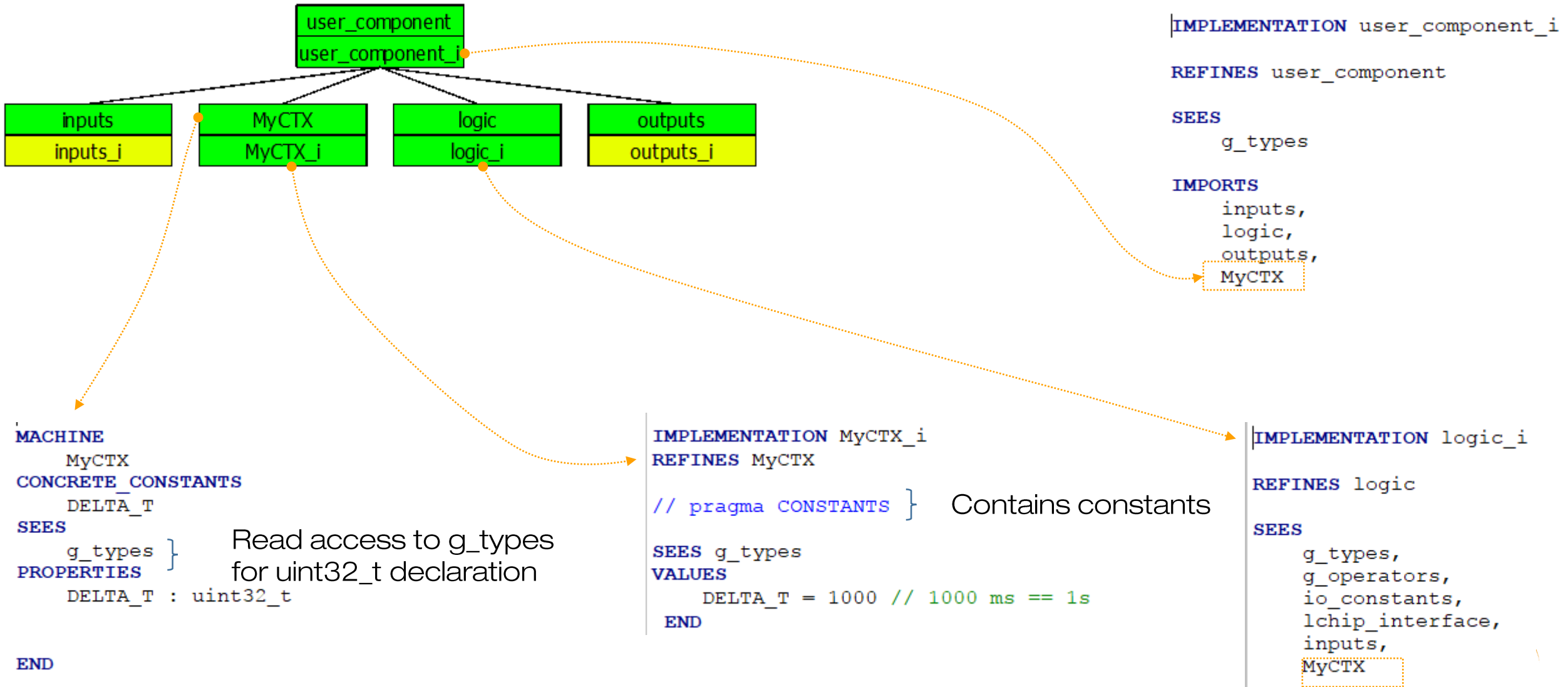
Example #2: clock

$O_1 = \text{not}(O_1)$ every 1 second

$O_2 = \text{not}(O_1)$

```
user_logic =  
BEGIN  
    VAR ltime_, s_tick_cycle, status_ IN  
        ltime_ : (ltime_ : uint32_t);  
        s_tick_cycle : (s_tick_cycle : uint32_t);  
        status_ : (status_ : uint32_t);  
  
    Current time — ltime_ <-- get_ms_tick;  
    s_tick_cycle := ltime_ / DELTA T; — User defined constant  
    status_ := s_tick_cycle mod 2;  
  
    O2 := IO_OFF;  
    IF status_ = 0 THEN  
        O2 := IO_ON  
    END;  
  
    IF O2 = IO_ON THEN  
        O1 := IO_OFF  
    ELSE  
        O1 := IO_ON  
    END  
END  
END;
```

Example #2: clock



Example #2: clock

Your turn:

- **Program 2 clocks with different cycle time** (one on O_1 , the other on O_2)
 $O_1 = \text{not}(O_1)$ every xxx milli-seconds
 $O_2 = \text{not}(O_2)$ every yyy milli-seconds

- **Do not change the status of the outputs too often (< 50 ms) or you will kill the relays !**
- Model, prove, compile, upload, reset your device
- Check with your buttons that the function is correctly implemented
- **Program clock which freezes when all 3 inputs are ON**

Wrap-up

Proof of Concept of the technology

Board SK₀ will be available for education July 2018

2 configurations:

- epoxy, switches, USB cable
- bare board

IDE updated 4 times / year

Pedagogical kit containing:

- a reference manual,
- a user manual
- models and programs
- examples including other computers

<http://www.clearsy.com/en/our-tools/clearsy-safety-platform/>

